# Attribute－Based Encryption 

陳榮傑
交通大學資工系
Cryptanalysis Lab
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## Public Key Encryption



Doctor encrypts EMR under Bob's public key
Doctor

Bob


Bob decrypts EMR under his private key

## Limitations

## $\square$ Bob is the single and known recipient of data

- Unknown recipient?
- Many recipients?
- More may join system later?


## Attribute-Based Encryption [swos]

Flexible data sharing:


## Idea

Ciphertexts: associated with access formulae


Private Keys: associated with attributes


Decryption:


## Example 1: Job Posting

Encrypt a job posting


## Example 2: Police Department

An informant encrypts a message for anyone in the internal affairs office or anyone who is undercover and in the central office.


## Example 3: Technology Company


hire date < 2002

## Example 4: Johns Hopkins Hospital



## Example 5: EMR (Electronic medical record)



## Example 5: EMR (Electronic medical record)



## Example 5: EMR



## Example 5: EMR



## Example 5: EMR



## Avoid Collusion Attacks

Keys must be personalized


## Key Personalization


"Vs"
Choose random r for each user's all attributes


## [BSW07,LW11] CIPHER-POLICY ABE

## Cipher-policy ABE

- Secret keys are labeled with a set of attributes
- Ciphertext is associated with access structure that control which user is able to decrypt the ciphertext.


## Setup

- Bilinear map: e
$-\mathrm{e}: \mathrm{G}_{1} \times \mathrm{G}_{1}->\mathrm{G}_{2}$
$-G_{1}$ has prime order $p$
-g is a generator of $\mathrm{G}_{1}$


## Setup

- $\mathrm{U}=\left\{\mathrm{a}_{1}=\right.$ child, $\left.\mathrm{a}_{2}=<120 \mathrm{~cm}, \ldots, \mathrm{a}_{\mathrm{n}}\right\}$
- $U$ is the set of all attributes
- $H: U->G_{1}$



## Setup

- MK(master key): used to produce user's secret key
- Choose $\alpha, \beta \in Z_{p}$
$-M K=\left(\beta, g^{\alpha}\right)$
- PK(public key): used to produce ciphertext
- PK=(g, $\left.g^{\beta}, e(g, g)^{\alpha}\right)$


## Encryption

- Encrypt(M(plaintext), T(access tree), PK) Choose a polynomial $q_{x}$ for each node: $q_{1}, q_{2}, q_{3}, \ldots, q_{8}$.

$$
\begin{aligned}
\text { degree }\left(q_{x}\right) & =K(x)-1 \\
\text { degree }\left(q_{1}\right) & =0 \\
\text { degree }\left(q_{2}\right) & =1 \\
\text { degree }\left(q_{3}\right) & =1 \\
\text { degree }\left(q_{4}\right) & =0
\end{aligned}
$$

- Encrypt(M(plaintext), T(access tree), PK)

Choose a polynomial $q_{x}$ for each node: $q_{1}, q_{2}, q_{3}, \ldots, q_{8}$.

```
degree(qx) = K(x) - 1
    degree(q}\mp@subsup{q}{1}{})=
    degree(q}\mp@subsup{q}{2}{})=
    degree(qu})=
    degree(qu) = 0
    degree (q8) =0
```


## Encryption



## Encryption

- Output
$-\mathrm{T}, \mathrm{Me}(\mathrm{g}, \mathrm{g})^{\alpha \mathrm{s}}, \mathrm{C}=\mathrm{g} \beta \mathrm{s}$
$-\mathrm{C} 4=\mathrm{g}^{\mathrm{q}_{4}(0)}$
C8
$-\mathrm{C4}{ }^{\prime}=\mathrm{H}(\text { child })^{q_{4}(0)}$ C8,



## Key Generation

KeyGen( $\gamma=$ \{ "child", "student", "<20" \}, MK )

- Choose $r \in Z_{p}$
- Choose $r_{\text {child }}, r_{\text {student }}, r_{<20} \in Z_{p}$
- Output
- $\mathrm{D}=\mathrm{g}^{(\alpha+r) / \beta}$
- Dchild $=\mathrm{g}^{r} \times H$ (child) ${ }^{r}$ child

Dstudent
D<20

- $\mathrm{D}^{\prime}$ child $=\mathrm{g}^{r_{\text {child }}}$

D'student
D'<20

## Decryption

- Cipher text C
$-\mathrm{T}, \mathrm{Me}(\mathrm{g}, \mathrm{g})^{\alpha \mathrm{s}}, \mathrm{C}=\mathrm{g}^{\beta \mathrm{s}}$
$-\mathrm{C} 4=\mathrm{g}_{4}(0)$

- $\mathrm{C4}^{\prime}=\mathrm{H}(\text { child })^{q_{4}(0)}$


C8'

- Private Key
- $\mathrm{D}=\mathrm{g}^{(\alpha+r) / \beta}$
- Dchild $=\mathrm{g}^{r} \times H(\text { child })^{r_{c h i l d}}$ Dstudent
D<20
- $\mathrm{D}^{\prime}$ child $=\mathrm{g}^{r_{\text {child }}}$

D'student
$D^{\prime}<20$


- $\frac{\mathrm{e}\left({\left.\text { Dstudent }, C_{6}\right)}_{\mathrm{e}\left(\mathrm{D}_{\text {student }}{ }^{\prime} \mathrm{C}^{\prime} 6\right)}\right.}{}=\mathrm{e}(\mathrm{g}, \mathrm{g})^{\mathrm{rq}_{6}(0)}$

$$
\begin{aligned}
& =\frac{e\left(D_{i}, C_{x}\right)}{e\left(D_{i}^{\prime}, C_{x}^{\prime}\right)} \\
& =\frac{e\left(g^{r} \cdot H(i)^{r_{i}}, h^{q_{x}(0)}\right)}{e\left(g^{r_{i}}, H(i)^{q_{x}(0)}\right)} \\
& =e(g, g)^{r q_{x}(0)} .
\end{aligned}
$$

- $e(g, g)^{r q_{1}(0)}=e(g, g)^{r s}$

$$
\begin{aligned}
& \operatorname{Me}(g, g)^{\alpha s} / e(C, D) \\
& \left.=\operatorname{Me}(g, g)^{\alpha s} / e^{\beta s}, g^{(\alpha+r) / \beta}\right) \\
& =\operatorname{Me}(g, g)^{-r s}
\end{aligned}
$$

$\mathrm{Me}(\mathrm{g}, \mathrm{g})^{-\mathrm{rs}} \cdot \mathrm{e}(\mathrm{g}, \mathrm{g})^{\mathrm{rs}}=\mathrm{M}$

## Our implementation (Linear Secret Sharing Scheme)

$$
\begin{aligned}
& \left(\begin{array}{ccc}
1 & 1 & 0 \\
0 & -1 & 0 \\
1 & 0 & 1 \\
1 & 0 & 2 \\
1 & 0 & 3
\end{array}\right) \cdot\left(\begin{array}{l}
s \\
r_{1} \\
r_{2}
\end{array}\right)= \\
& \left(\begin{array}{c}
s+r_{1} \\
-r_{1} \\
s+r_{2} \\
s+2 r_{2} \\
s+3 r_{2}
\end{array}\right) \begin{array}{l}
(\text { for } 4) \\
\left(\begin{array}{l}
\text { (for 5) }
\end{array}\right. \\
(\text { for } 6) \\
\text { (for } 7) \\
\text { (for } 8)
\end{array}
\end{aligned}
$$




## 病歷首頁單



Continuity of Care Document：
XML－based standard

Access Formula
"_:Pathology" OR ("_:VS" AND "_:Surgery")


## encrypt!

Ciphers


Ciphers

$$
\begin{array}{|l|l|}
\hline 25 . a e s 128 & \text { policy? } \\
\hline
\end{array}
$$



## 病歷首頁單



- 重大傷病：（文字敘述）
- 過敏史－藥物過敏：（文字敘述）
- 過敏史－食物過敏：（文字敘述）
- 過敏史－環境過敏：（文字敘述）
- 藥物不良反應（ADR）：（文字敘述）
- 旅遊史：（文字敘述）
- 傳染病史：（文字敘述）
- 遺傳病史：（文字敘述）
- 其他個人重要病史：（文字敘述）

- Questions?
- Thank you

